Getting started

Coq is a Proof Assistant for a Logical Framework known as the Calculus of Inductive Constructions. It allows the interactive construction of formal proofs, and also the manipulation of functional programs consistently with their specifications. It runs as a computer program on many architectures. It is available with a variety of user interfaces. The present document does not attempt to present a comprehensive view of all the possibilities of Coq, but rather to present in the most elementary manner a tutorial on the basic specification language, called Gallina, in which formal axiomatisations may be developed, and on the main proof tools. For more advanced information, the reader could refer to the Coq Reference Manual or the Coq’Art, a book by Y. Bertot and P. Castéran on practical uses of the Coq system.

Coq can be used from a standard teletype-like shell window but preferably through the graphical user interface CoqIde.

Instructions on installation procedures, as well as more comprehensive documentation, may be found in the standard distribution of Coq, which may be obtained from Coq web site [https://coq.inria.fr/](https://coq.inria.fr/).

In the following, we assume that Coq is called from a standard teletype-like shell window. All examples preceded by the prompting sequence Coq < represent user input, terminated by a period.

The following lines usually show Coq’s answer as it appears on the users screen. When used from a graphical user interface such as CoqIde, the prompt is not displayed: user input is given in one window and Coq’s answers are displayed in a different window.

The sequence of such examples is a valid Coq session, unless otherwise specified. This version of the tutorial has been prepared on a PC workstation running Linux. The standard invocation of Coq delivers a message such as:

```
unix:~> coqtop
Welcome to Coq 8.2 (January 2009)
Coq <
```

The first line gives a banner stating the precise version of Coq used. You should always return this banner when you report an anomaly to our bug-tracking system [https://coq.inria.fr/bugs/](https://coq.inria.fr/bugs/).

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1 Alternative graphical interfaces exist: Proof General and Pcoq.
Chapter 1

Basic Predicate Calculus

1.1 An overview of the specification language Gallina

A formal development in Gallina consists in a sequence of declarations and definitions. You may also send Coq commands which are not really part of the formal development, but correspond to information requests, or service routine invocations. For instance, the command:

Coq < Quit.

terminates the current session.

1.1.1 Declarations

A declaration associates a name with a specification. A name corresponds roughly to an identifier in a programming language, i.e. to a string of letters, digits, and a few ASCII symbols like underscore (\_) and prime ('), starting with a letter. We use case distinction, so that the names A and a are distinct. Certain strings are reserved as key-words of Coq, and thus are forbidden as user identifiers.

A specification is a formal expression which classifies the notion which is being declared. There are basically three kinds of specifications: logical propositions, mathematical collections, and abstract types. They are classified by the three basic sorts of the system, called respectively Prop, Set, and Type, which are themselves atomic abstract types.

Every valid expression \( e \) in Gallina is associated with a specification, itself a valid expression, called its type \( \tau(E) \). We write \( e : \tau(E) \) for the judgment that \( e \) is of type \( E \). You may request Coq to return to you the type of a valid expression by using the command Check:

Coq < Check 0.
0
  : nat
Thus we know that the identifier \( O \) (the name ‘O’, not to be confused with the numeral ‘0’ which is not a proper identifier!) is known in the current context, and that its type is the specification \( \text{nat} \). This specification is itself classified as a mathematical collection, as we may readily check:

Coq < Check nat.
\[ \text{nat} : \text{Set} \]

The specification \( \text{Set} \) is an abstract type, one of the basic sorts of the Gallina language, whereas the notions \( \text{nat} \) and \( O \) are notions which are defined in the arithmetic prelude, automatically loaded when running the COQ system.

We start by introducing a so-called section name. The role of sections is to structure the modelisation by limiting the scope of parameters, hypotheses and definitions. It will also give a convenient way to reset part of the development.

Coq < Section Declaration.

With what we already know, we may now enter in the system a declaration, corresponding to the informal mathematics *let n be a natural number*.

Coq < Variable n : nat.
\[ n \text{ is declared} \]

If we want to translate a more precise statement, such as *let n be a positive natural number*, we have to add another declaration, which will declare explicitly the hypothesis \( \text{Pos}_n \), with specification the proper logical proposition:

Coq < Hypothesis Pos_n : (gt n 0).
\[ \text{Pos}_n \text{ is declared} \]

Indeed we may check that the relation \( \text{gt} \) is known with the right type in the current context:

Coq < Check gt.
\[ \text{gt} : \text{nat} \to \text{nat} \to \text{Prop} \]

which tells us that \( \text{gt} \) is a function expecting two arguments of type \( \text{nat} \) in order to build a logical proposition. What happens here is similar to what we are used to in a functional programming language: we may compose the (specification) type \( \text{nat} \) with the (abstract) type \( \text{Prop} \) of logical propositions through the arrow function constructor, in order to get a functional type \( \text{nat}\to\text{Prop} \):

Coq < Check (nat -> Prop).
\[ \text{nat} \to \text{Prop} \]
\[ : \text{Type} \]
which may be composed one more times with nat in order to obtain the type nat->nat->Prop of binary relations over natural numbers. Actually the type nat->nat->Prop is an abbreviation for nat->(nat->Prop).

Functional notions may be composed in the usual way. An expression f of type A → B may be applied to an expression e of type A in order to form the expression \( f(e) \) of type B. Here we get that the expression \( \text{gt } n \) is well-formed of type nat->Prop, and thus that the expression \( \text{gt } n \text{0} \), which abbreviates \( ((\text{gt } n) \text{0}) \), is a well-formed proposition.

Coq < Check gt n 0.
\( n > 0 \) : Prop

1.1.2 Definitions

The initial prelude contains a few arithmetic definitions: nat is defined as a mathematical collection (type Set), constants O, S, plus, are defined as objects of types respectively nat, nat->nat, and nat->nat->nat. You may introduce new definitions, which link a name to a well-typed value. For instance, we may introduce the constant one as being defined to be equal to the successor of zero:

Coq < Definition one := (S O).

We may optionally indicate the required type:

Coq < Definition two : nat := S one.

Actually COQ allows several possible syntaxes:

Coq < Definition three := S two : nat.

Here is a way to define the doubling function, which expects an argument m of type nat in order to build its result as \( \text{plus } m \text{m} \):  

Coq < Definition double (m:nat) := plus m m.

double is defined

This introduces the constant double defined as the expression fun m:nat => plus m m. The abstraction introduced by fun is explained as follows. The expression fun x:A => e is well formed of type A→B in a context whenever the expression e is well-formed of type B in the given context to which we add the declaration that x is of type A. Here x is a bound, or dummy variable in the expression fun x:A => e. For instance we could as well have defined double as fun n:nat => (plus n n).

Bound (local) variables and free (global) variables may be mixed. For instance, we may define the function which adds the constant n to its argument as
Coq < Definition add_n (m:nat) := plus m n.
add_n is defined

However, note that here we may not rename the formal argument \( m \) into \( n \) without capturing the free occurrence of \( n \), and thus changing the meaning of the defined notion.

Binding operations are well known for instance in logic, where they are called quantifiers. Thus we may universally quantify a proposition such as \( m > 0 \) in order to get a universal proposition \( \forall m \cdot m > 0 \). Indeed this operator is available in Coq, with the following syntax: \( \forall m : \text{nat}, \, \text{gt m 0} \). Similarly to the case of the functional abstraction binding, we are obliged to declare explicitly the type of the quantified variable. We check:

Coq < Check (forall m:nat, gt m 0).
forall m : nat, m > 0 : Prop

We may revert to the clean state of our initial session using the Coq Reset command:

Coq < Reset Initial.

1.2 Introduction to the proof engine: Minimal Logic

In the following, we are going to consider various propositions, built from atomic propositions \( A, B, C \). This may be done easily, by introducing these atoms as global variables declared of type Prop. It is easy to declare several names with the same specification:

Coq < Section Minimal_Logic.
Coq < Variables A B C : Prop.
A is declared
B is declared
C is declared

We shall consider simple implications, such as \( A \to B \), read as “\( A \implies B \)”. Remark that we overload the arrow symbol, which has been used above as the functionality type constructor, and which may be used as well as propositional connective:

Coq < Check (A -> B).
A -> B : Prop

Let us now embark on a simple proof. We want to prove the easy tautology \(( (A \to (B \to C)) \to (A \to B) \to (A \to C) \)\). We enter the proof engine by the command Goal, followed by the conjecture we want to verify:

Coq < Goal (A -> B -> C) -> (A -> B) -> A -> C.
1 subgoal
1.2. INTRODUCTION TO THE PROOF ENGINE: MINIMAL LOGIC

\[ A, B, C : \text{Prop} \]
\[ \begin{align*}
\text{A} \rightarrow \text{B} \rightarrow \text{C} & \rightarrow (\text{A} \rightarrow \text{B}) \rightarrow \text{A} \rightarrow \text{C} \\
\end{align*} \]

The system displays the current goal below a double line, local hypotheses (there are none initially) being displayed above the line. We call the combination of local hypotheses with a goal a judgment. We are now in an inner loop of the system, in proof mode. New commands are available in this mode, such as tactics, which are proof combining primitives. A tactic operates on the current goal by attempting to construct a proof of the corresponding judgment, possibly from proofs of some hypothetical judgments, which are then added to the current list of conjectured judgments. For instance, the intro tactic is applicable to any judgment whose goal is an implication, by moving the proposition to the left of the application to the list of local hypotheses:

Coq < intro H.
1 subgoal

\[ A, B, C : \text{Prop} \]
\[ H : A \rightarrow B \rightarrow C \]
\[ \begin{align*}
\text{A} \rightarrow B & \rightarrow \text{A} \rightarrow \text{C} \\
\end{align*} \]

Several introductions may be done in one step:

Coq < intros H' HA.
1 subgoal

\[ A, B, C : \text{Prop} \]
\[ H : A \rightarrow B \rightarrow C \]
\[ H' : A \rightarrow B \]
\[ HA : A \]
\[ \begin{align*}
\text{C} \\
\end{align*} \]

We notice that \( \text{C} \), the current goal, may be obtained from hypothesis \( H \), provided the truth of \( A \) and \( B \) are established. The tactic apply implements this piece of reasoning:

Coq < apply H.
2 subgoals

\[ A, B, C : \text{Prop} \]
\[ H : A \rightarrow B \rightarrow C \]
\[ H' : A \rightarrow B \]
\[ HA : A \]
\[ \begin{align*}
\text{A} \\
\text{subgoal 2 is:} \\
\text{B} \\
\end{align*} \]
We are now in the situation where we have two judgments as conjectures that remain to be proved. Only the first is listed in full, for the others the system displays only the corresponding subgoal, without its local hypotheses list. Remark that \texttt{apply} has kept the local hypotheses of its father judgment, which are still available for the judgments it generated.

In order to solve the current goal, we just have to notice that it is exactly available as hypothesis \texttt{HA}:

\begin{verbatim}
Coq < exact HA.
1 subgoal

A, B, C : Prop
H : A -> B -> C
H' : A -> B
HA : A

B

Now H' applies:

Coq < apply H'.
1 subgoal

A, B, C : Prop
H : A -> B -> C
H' : A -> B
HA : A

A

And we may now conclude the proof as before, with \texttt{exact HA}. Actually, we may not bother with the name \texttt{HA}, and just state that the current goal is solvable from the current local assumptions:

\begin{verbatim}
Coq < assumption.
No more subgoals.
\end{verbatim}

The proof is now finished. We may either discard it, by using the command \texttt{Abort} which returns to the standard \texttt{COQ} toplevel loop without further ado, or else save it as a lemma in the current context, under name say \texttt{trivial_lemma}:

\begin{verbatim}
Coq < Save trivial_lemma.
intro H.
(intros H' HA).
(apply H).
exact HA.
(apply H').
assumption.
Save trivial_lemma.
trivial_lemma is defined
\end{verbatim}
As a comment, the system shows the proof script listing all tactic commands used in the proof.

Let us redo the same proof with a few variations. First of all we may name the initial goal as a conjectured lemma:

Coq < Lemma distr_impl : (A -> B -> C) -> (A -> B) -> A -> C.
1 subgoal
A, B, C : Prop
(\(A \rightarrow B \rightarrow C\)) \(\rightarrow\) (\(A \rightarrow B\)) \(\rightarrow\) A \(\rightarrow\) C

Next, we may omit the names of local assumptions created by the introduction tactics, they can be automatically created by the proof engine as new non-clashing names.

Coq < intros.
1 subgoal
A, B, C : Prop
H : A \rightarrow B \rightarrow C
H0 : A \rightarrow B
H1 : A
C

The intros tactic, with no arguments, effects as many individual applications of intro as is legal.

Then, we may compose several tactics together in sequence, or in parallel, through tacticals, that is tactic combinators. The main constructions are the following:

- \(T_1; T_2\) (read \(T_1\) then \(T_2\)) applies tactic \(T_1\) to the current goal, and then tactic \(T_2\) to all the subgoals generated by \(T_1\).
- \(T; [T_1 | T_2 | ... | T_n]\) applies tactic \(T\) to the current goal, and then tactic \(T_1\) to the first newly generated subgoal, ...., \(T_n\) to the nth.

We may thus complete the proof of distr_impl with one composite tactic:

Coq < apply H; [ assumption | apply H0; assumption ].
No more subgoals.

Let us now save lemma distr_impl:

Coq < Save.
(intros **).
(appy H; [ assumption | apply H0; assumption ]).
Qed.

distr_impl is defined
Here \texttt{Save} needs no argument, since we gave the name \texttt{distr_impl} in advance; it is however possible to override the given name by giving a different argument to command \texttt{Save}.

Actually, such an easy combination of tactics \texttt{intro}, \texttt{apply} and \texttt{assumption} may be found completely automatically by an automatic tactic, called \texttt{auto}, without user guidance:

\begin{verbatim}
Coq < Lemma distr_imp : (A -> B -> C) -> (A -> B) -> A -> C.
1 subgoal
A, B, C : Prop
============================
(A -> B -> C) -> (A -> B) -> A -> C
Coq < auto.
No more subgoals.
\end{verbatim}

This time, we do not save the proof, we just discard it with the \texttt{Abort} command:

\begin{verbatim}
Coq < Abort.
\end{verbatim}

At any point during a proof, we may use \texttt{Abort} to exit the proof mode and go back to Coq's main loop. We may also use \texttt{Restart} to restart from scratch the proof of the same lemma. We may also use \texttt{Undo} to backtrack one step, and more generally \texttt{Undo n} to backtrack n steps.

We end this section by showing a useful command, \texttt{Inspect n.}, which inspects the global COQ environment, showing the last n declared notions:

\begin{verbatim}
Coq < Inspect 3.
*** [C : Prop]
 trivial_lemma : (A -> B -> C) -> (A -> B) -> A -> C
distr_impl : (A -> B -> C) -> (A -> B) -> A -> C
\end{verbatim}

The declarations, whether global parameters or axioms, are shown preceded by 
***; definitions and lemmas are stated with their specification, but their value (or proof-term) is omitted.

1.3 Propositional Calculus

1.3.1 Conjunction

We have seen how \texttt{intro} and \texttt{apply} tactics could be combined in order to prove implicational statements. More generally, COQ favors a style of reasoning, called \textit{Natural Deduction}, which decomposes reasoning into so called \textit{introduction rules}, which tell how to prove a goal whose main operator is a given propositional connective, and \textit{elimination rules}, which tell how to use an hypothesis whose main operator is the propositional connective. Let us show how to use these ideas for the propositional connectives $\land$ and $\lor$.
Coq < Lemma and_commutative : A \ B -> B \ A.
1 subgoal

A, B : Prop

We make use of the conjunctive hypothesis \( H \) with the elim tactic, which breaks it into its components:

Coq < elim H.
1 subgoal

A, B : Prop
H : A \ B

We now use the conjunction introduction tactic split, which splits the conjunctive goal into the two subgoals:

Coq < split.
2 subgoals

A, B : Prop
H : A \ B
H0 : A
H1 : B

and the proof is now trivial. Indeed, the whole proof is obtainable as follows:

Coq < Restart.
1 subgoal

A, B : Prop

No more subgoals.
1.3.2 Disjunction

In a similar fashion, let us consider disjunction:

Coq < Lemma or_commutative : A \ B -> B \ A.
  1 subgoal
    A, B, C : Prop
    =============================
    A \ B -> B \ A
  Coq < intro H; elim H.
  2 subgoals
    A, B, C : Prop
    H : A \ B
    =============================
    A -> B \ A
  subgoal 2 is:
  B -> B \ A

Let us prove the first subgoal in detail. We use intro in order to be left to prove B\A from H:

Coq < intro HA.
  2 subgoals
    A, B, C : Prop
    H : A \ B
    HA : A
    =============================
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\[ B \lor A \]

subgoal 2 is:
\[ B \rightarrow B \lor A \]

Here the hypothesis \( H \) is not needed anymore. We could choose to actually erase it with the tactic `clear`; in this simple proof it does not really matter, but in bigger proof developments it is useful to clear away unnecessary hypotheses which may clutter your screen.

```coq
Coq < clear H.
2 subgoals

A, B, C : Prop
HA : A

=======================================
B \lor A
subgoal 2 is:
B \rightarrow B \lor A
```

The disjunction connective has two introduction rules, since \( P \lor Q \) may be obtained from \( P \) or from \( Q \); the two corresponding proof constructors are called respectively `or_introl` and `or_intror`; they are applied to the current goal by tactics `left` and `right` respectively. For instance:

```coq
Coq < right.
2 subgoals

A, B, C : Prop
HA : A

=======================================
A
subgoal 2 is:
B \rightarrow B \lor A
```

```coq
Coq < trivial.
1 subgoal

A, B, C : Prop
H : A \lor B

=======================================
B \rightarrow B \lor A
```

The tactic `trivial` works like `auto` with the hints database, but it only tries those tactics that can solve the goal in one step.

As before, all these tedious elementary steps may be performed automatically, as shown for the second symmetric case:

```coq
Coq < auto.
No more subgoals.
```
However, auto alone does not succeed in proving the full lemma, because it does not try any elimination step. It is a bit disappointing that auto is not able to prove automatically such a simple tautology. The reason is that we want to keep auto efficient, so that it is always effective to use.

1.3.3 Tauto

A complete tactic for propositional tautologies is indeed available in COQ as the tauto tactic.

Coq < Restart.
1 subgoal

\[ A, B, C : \text{Prop} \]
\[ \text{============================} \]
\[ A \lor B \rightarrow B \lor A \]

Coq < tauto.
No more subgoals.

Coq < Qed.
tauto.
Qed.
or_commutative is defined

It is possible to inspect the actual proof tree constructed by tauto, using a standard command of the system, which prints the value of any notion currently defined in the context:

Coq < Print or_commutative.

\[
\text{or_commutative = } \\
\text{fun H : A \lor B =>} \\
\text{or_ind (fun H0 : A => or_intror H0)} \\
\text{(fun H0 : B => or_introl H0) H} \\
\text{: A \lor B \rightarrow B \lor A}
\]

It is not easy to understand the notation for proof terms without a few explanations. The fun prefix, such as \text{fun H:A\lor B =>}, corresponds to intro H, whereas a subterm such as \text{(or_intror B H0)} corresponds to the sequence of tactics apply or_intror; exact H0. The generic combinator or_intror needs to be instantiated by the two properties B and A. Because A can be deduced from the type of H0, only B is printed. The two instantiations are effected automatically by the tactic apply when pattern-matching a goal. The specialist will of course recognize our proof term as a \(\lambda\)-term, used as notation for the natural deduction proof term through the Curry-Howard isomorphism. The naive user of COQ may safely ignore these formal details.

Let us exercise the tauto tactic on a more complex example:

Coq < Lemma distr_and : A \rightarrow B \lor C \rightarrow (A \rightarrow B) \lor (A \rightarrow C).
1 subgoal
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\[ A, B, C : \text{Prop} \]
\[
\begin{array}{c}
\text{=====}
\end{array}
\]
\[ A \rightarrow B \land C \rightarrow (A \rightarrow B) \land (A \rightarrow C) \]

Coq < tauto.
No more subgoals.
Coq < Qed.
tauto.
Qed.
distr_and is defined

1.3.4 Classical reasoning

The tactic tauto always comes back with an answer. Here is an example where it fails:

Coq < Lemma Peirce : ((A \rightarrow B) \rightarrow A) -> A.
1 subgoal

\[ A, B, C : \text{Prop} \]
\[
\begin{array}{c}
\text{=====}
\end{array}
\]
\[ ((A \rightarrow B) \rightarrow A) \rightarrow A \]

Coq < try tauto.
1 subgoal

\[ A, B, C : \text{Prop} \]
\[
\begin{array}{c}
\text{=====}
\end{array}
\]
\[ ((A \rightarrow B) \rightarrow A) \rightarrow A \]

Note the use of the Try tactical, which does nothing if its tactic argument fails.

This may come as a surprise to someone familiar with classical reasoning. Peirce’s lemma is true in Boolean logic, i.e. it evaluates to true for every truth-assignment to \( A \) and \( B \). Indeed the double negation of Peirce’s law may be proved in Coq using tauto:

Coq < Abort.
Coq < Lemma NNPeirce : \~ \~ ((A \rightarrow B) \rightarrow A) \rightarrow A).
1 subgoal

\[ A, B, C : \text{Prop} \]
\[
\begin{array}{c}
\text{=====}
\end{array}
\]
\[ \sim \sim ((A \rightarrow B) \rightarrow A) \rightarrow A \]

Coq < tauto.
No more subgoals.
Coq < Qed.
tauto.
Qed.
NNPeirce is defined
In classical logic, the double negation of a proposition is equivalent to this proposition, but in the constructive logic of COQ this is not so. If you want to use classical logic in COQ, you have to import explicitly the Classical module, which will declare the axiom \texttt{classic} of excluded middle, and classical tautologies such as de Morgan’s laws. The \texttt{Require} command is used to import a module from COQ’s library:

\begin{verbatim}
Coq < Require Import Classical.

Coq < Check NNPP.
NNPP
   : forall p : Prop, ~ ~ p -> p

and it is now easy (although admittedly not the most direct way) to prove a classical law such as Peirce’s:

Coq < Lemma Peirce : ((A -> B) -> A) -> A.
1 subgoal

A, B, C : Prop

===========================
((A -> B) -> A) -> A

Coq < apply NNPP; tauto.
No more subgoals.

Coq < Qed.
(apply NNPP; tauto).
Qed.
Peirce is defined
\end{verbatim}

Here is one more example of propositional reasoning, in the shape of a Scottish puzzle. A private club has the following rules:

1. Every non-scottish member wears red socks
2. Every member wears a kilt or doesn’t wear red socks
3. The married members don’t go out on Sunday
4. A member goes out on Sunday if and only if he is Scottish
5. Every member who wears a kilt is Scottish and married
6. Every scottish member wears a kilt

Now, we show that these rules are so strict that no one can be accepted.

\begin{verbatim}
Coq < Section club.

Scottish is declared
RedSocks is declared
WearKilt is declared
\end{verbatim}
Married is declared
GoOutSunday is declared

Coq < Hypothesis rule1 : ~ Scottish -> RedSocks.
rule1 is declared

Coq < Hypothesis rule2 : WearKilt \/ ~ RedSocks.
rule2 is declared

Coq < Hypothesis rule3 : Married -> ~ GoOutSunday.
rule3 is declared

Coq < Hypothesis rule4 : GoOutSunday <-> Scottish.
rule4 is declared

Coq < Hypothesis rule5 : WearKilt -> Scottish /\ Married.
rule5 is declared

Coq < Hypothesis rule6 : Scottish -> WearKilt.
rule6 is declared

Coq < Lemma NoMember : False.
1 subgoal

A, B, C, Scottish, RedSocks, WearKilt, Married,
GoOutSunday : Prop
rule1 : ~ Scottish -> RedSocks
rule2 : WearKilt \/ ~ RedSocks
rule3 : Married -> ~ GoOutSunday
rule4 : GoOutSunday <-> Scottish
rule5 : WearKilt -> Scottish /\ Married
rule6 : Scottish -> WearKilt

============================
False

Coq < tauto.
No more subgoals.

Coq < Qed.
tauto.
Qed.
NoMember is defined

At that point NoMember is a proof of the absurdity depending on hypotheses. We may end the section, in that case, the variables and hypotheses will be discharged, and the type of NoMember will be generalised.

Coq < End club.

Coq < Check NoMember.
NoMember : forall
   Scottish RedSocks WearKilt Married
   GoOutSunday : Prop,
   (~ Scottish -> RedSocks) ->
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\[
\begin{align*}
\text{WearKilt} \lor \neg \text{RedSocks} & \rightarrow \\
(\text{Married} \rightarrow \neg \text{GoOutSunday}) & \rightarrow \\
\text{GoOutSunday} & \leftrightarrow \text{Scottish} \rightarrow \\
(\text{WearKilt} \rightarrow \neg \text{Married}) & \rightarrow \\
(\text{Scottish} \rightarrow \text{WearKilt}) & \rightarrow \text{False}
\end{align*}
\]

1.4 Predicate Calculus

Let us now move into predicate logic, and first of all into first-order predicate calculus. The essence of predicate calculus is that to try to prove theorems in the most abstract possible way, without using the definitions of the mathematical notions, but by formal manipulations of uninterpreted function and predicate symbols.

1.4.1 Sections and signatures

Usually one works in some domain of discourse, over which range the individual variables and function symbols. In COQ we speak in a language with a rich variety of types, so one may mix several domains of discourse, in our multi-sorted language. For the moment, we just do a few exercises, over a domain of discourse \( D \) axiomatised as a Set, and we consider two predicate symbols \( P \) and \( R \) over \( D \), of arities respectively 1 and 2. Such abstract entities may be entered in the context as global variables. But we must be careful about the pollution of our global environment by such declarations. For instance, we have already polluted our COQ session by declaring the variables \( n, \text{Pos}_n, A, B, \) and \( C \).

Coq < Reset Initial.

We shall now declare a new Section, which will allow us to define notions local to a well-delimited scope. We start by assuming a domain of discourse \( D \), and a binary relation \( R \) over \( D \):

Coq < Section Predicate_calculus.

Coq < Variable D : Set.

D is declared

Coq < Variable R : D -> D -> Prop.

R is declared

As a simple example of predicate calculus reasoning, let us assume that relation \( R \) is symmetric and transitive, and let us show that \( R \) is reflexive in any point \( x \) which has an \( R \) successor. Since we do not want to make the assumptions about \( R \) global axioms of a theory, but rather local hypotheses to a theorem, we open a specific section to this effect.

Coq < Section R_sym_trans.

Coq < Hypothesis R_symmetric : \( \forall x y : D, R x y \rightarrow R y x \).

R_symmetric is declared

Coq < Hypothesis R_transitive : \( \forall x y z : D, R x y \rightarrow R y z \rightarrow R x z \).

R_transitive is declared
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Remark the syntax \( \forall x : D \), which stands for universal quantification \( \forall x : D \).

1.4.2 Existential quantification

We now state our lemma, and enter proof mode.

Coq < Lemma refl_if : forall x : D, (exists y, R x y) -> R x x.
1 subgoal

\[
\begin{align*}
D & : \text{Set} \\
R & : D \rightarrow D \rightarrow \text{Prop} \\
R_{\text{symmetric}} & : \forall x y : D, R x y \rightarrow R y x \\
R_{\text{transitive}} & : \forall x y z : D, R x y \rightarrow R y z \rightarrow R x z \\
\end{align*}
\]

\[
\begin{array}{c}
\text{forall } x : D, (\exists y : D, R x y) \rightarrow R x x \\
\end{array}
\]

Remark that the hypotheses which are local to the currently opened sections are listed as local hypotheses to the current goals. The rationale is that these hypotheses are going to be discharged, as we shall see, when we shall close the corresponding sections.

Note the functional syntax for existential quantification. The existential quantifier is built from the operator \( \text{ex} \), which expects a predicate as argument:

Coq < Check ex.

ex

: forall A : Type, (A -> Prop) -> Prop

and the notation \( \exists x : D, P x \) is just concrete syntax for the expression \( \text{ex} D \ (\text{fun} x : D \Rightarrow P x) \). Existential quantification is handled in COQ in a similar fashion to the connectives /\ and /\ : it is introduced by the proof combinator \( \text{ex_intro} \), which is invoked by the specific tactic \( \text{Exists} \), and its elimination provides a witness \( a : D \) to \( P \), together with an assumption \( h : (P a) \) that indeed \( a \) verifies \( P \). Let us see how this works on this simple example.

Coq < intros x x_Rlinked.
1 subgoal

\[
\begin{align*}
D & : \text{Set} \\
R & : D \rightarrow D \rightarrow \text{Prop} \\
R_{\text{symmetric}} & : \forall x y : D, R x y \rightarrow R y x \\
R_{\text{transitive}} & : \forall x y z : D, R x y \rightarrow R y z \rightarrow R x z \\
x & : D \\
x_{\text{Rlinked}} & : \exists y : D, R x y \\
R x y & \\
\end{align*}
\]

Remark that \( \text{intros} \) treats universal quantification in the same way as the premises of implications. Renaming of bound variables occurs when it is needed; for instance, had we started with \( \text{intro } y \), we would have obtained the goal:
Coq < intro y.
1 subgoal

D : Set
R : D -> D -> Prop
R_symmetric : forall x y : D, R x y -> R y x
R_transitive : forall x y z : D, R x y -> R y z -> R x z
y : D

Let us now use the existential hypothesis x_Rlinked to exhibit an R-successor y of x. This is done in two steps, first with elim, then with intros

Coq < elim x_Rlinked.
1 subgoal

D : Set
R : D -> D -> Prop
R_symmetric : forall x y : D, R x y -> R y x
R_transitive : forall x y z : D, R x y -> R y z -> R x z
x : D
x_Rlinked : exists y : D, R x y

Now we want to use R_transitive. The apply tactic will know how to match x with x, and z with x, but needs help on how to instantiate y, which appear in the hypotheses of R_transitive, but not in its conclusion. We give the proper hint to apply in a with clause, as follows:

Coq < apply R_transitive with y.
2 subgoals

D : Set
R : D -> D -> Prop
R_symmetric : forall x y : D, R x y -> R y x
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\[
\begin{align*}
R_{\text{transitive}} : & \quad \forall x \ y \ z : D, \ R \ x \ y \ \rightarrow \ R \ y \ z \ \rightarrow \ R \ x \ z \\
x \ : D \\
x_{\text{Rlinked}} : & \quad \exists y : D, \ R \ x \ y \\
y \ : D \\
R_{xy} : & \quad R \ x \ y \\
\end{align*}
\]

\[\begin{array}{l}
\text{subgoal 2 is:} \\
R \ y \ x
\end{array}\]

The rest of the proof is routine:

Coq < assumption.
1 subgoal

\[
\begin{align*}
D : & \quad \text{Set} \\
R : & \quad D \rightarrow D \rightarrow \text{Prop} \\
R_{\text{symmetric}} : & \quad \forall x \ y : D, \ R \ x \ y \ \rightarrow \ R \ y \ x \\
R_{\text{transitive}} : & \quad \forall x \ y \ z : D, \ R \ x \ y \ \rightarrow \ R \ y \ z \ \rightarrow \ R \ x \ z \\
x \ : D \\
x_{\text{Rlinked}} : & \quad \exists y : D, \ R \ x \ y \\
y \ : D \\
R_{xy} : & \quad R \ x \ y \\
\end{align*}
\]

Coq < apply R_symmetric; assumption.
No more subgoals.
Coq < Qed.

Let us now close the current section.

Coq < End R_sym_trans.

All the local hypotheses have been discharged in the statement of \text{refl_if},
which now becomes a general theorem in the first-order language declared in sec-
tion \text{Predicate_calculus}. In this particular example, the use of section \text{R_sym_trans}
has not been really significant, since we could have instead stated theorem \text{refl_if}
in its general form, and done basically the same proof, obtaining \text{R_symmetric} and
\text{R_transitive} as local hypotheses by initial \text{intros} rather than as global hypothe-
ses in the context. But if we had pursued the theory by proving more theorems
about relation \( R \), we would have obtained all general statements at the closing of
the section, with minimal dependencies on the hypotheses of symmetry and transi-
tivity.

1.4.3 Paradoxes of classical predicate calculus

Let us illustrate this feature by pursuing our \text{Predicate_calculus} section with
an enrichment of our language: we declare a unary predicate \( P \) and a constant \( d \):
We shall now prove a well-known fact from first-order logic: a universal predicate is non-empty, or in other terms existential quantification follows from universal quantification.

Coq < Lemma weird : (forall x:D, P x) -> exists a, P a.
1 subgoal

D : Set
R : D -> D -> Prop
P : D -> Prop
d : D

UnivP : forall x : D, P x
exists a : D, P a

First of all, notice the pair of parentheses around \(\forall x : D, P x\) in the statement of lemma weird. If we had omitted them, COQ's parser would have interpreted the statement as a truly trivial fact, since we would postulate an \(x\) verifying \((P x)\). Here the situation is indeed more problematic. If we have some element in \(D\), we may apply UnivP to it and conclude, otherwise we are stuck. Indeed such an element \(d\) exists, but this is just by virtue of our new signature. This points out a subtle difference between standard predicate calculus and COQ. In standard first-order logic, the equivalent of lemma weird always holds, because such a rule is wired in the inference rules for quantifiers, the semantic justification being that the interpretation domain is assumed to be non-empty. Whereas in COQ, where types are not assumed to be systematically inhabited, lemma weird only holds in signatures which allow the explicit construction of an element in the domain of the predicate.

Let us conclude the proof, in order to show the use of the Exists tactic:

Coq < exists d; trivial.
No more subgoals.
Coq < Qed.
intro UnivP.
(exists d; trivial).
Qed.

weird is defined

Another fact which illustrates the sometimes disconcerting rules of classical predicate calculus is Smullyan’s drinkers’ paradox: “In any non-empty bar, there is a person such that if she drinks, then everyone drinks”. We modelize the bar by Set D, drinking by predicate P. We shall need classical reasoning. Instead of loading the Classical module as we did above, we just state the law of excluded middle as a local hypothesis schema at this point:

Coq < Hypothesis EM : forall A:Prop, A \or ~ A.
EM is declared
Coq < Lemma drinker : exists x:D, P x -> forall x:D, P x.
1 subgoal

D : Set
R : D -> D -> Prop
P : D -> Prop
d : D
EM : forall A : Prop, A \or ~ A
============================
exists x : D, P x -> forall x0 : D, P x0

The proof goes by cases on whether or not there is someone who does not drink. Such reasoning by cases proceeds by invoking the excluded middle principle, via elim of the proper instance of EM:

Coq < elim (EM (exists x, ~ P x)).
2 subgoals

D : Set
R : D -> D -> Prop
P : D -> Prop
d : D
EM : forall A : Prop, A \or ~ A
============================
(exists x : D, ~ P x) ->
exists x : D, P x -> forall x0 : D, P x0
subgoal 2 is:
~ (exists x : D, ~ P x) ->
exists x : D, P x -> forall x0 : D, P x0

We first look at the first case. Let Tom be the non-drinker:

Coq < intro Non_drinker; elim Non_drinker;
intro Tom; Tom_does_not_drink.
2 subgoals

D : Set
R : D -> D -> Prop
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\[ P : D \rightarrow Prop \]
\[ d : D \]
\[ EM : \forall A : Prop, A \lor \neg A \]
\[ Non\_drinker : \exists x : D, \neg P x \]
\[ Tom : D \]
\[ Tom\_does\_not\_drink : \neg P Tom \]

\[ \exists x : D, P x \rightarrow \forall x0 : D, P x0 \]

subgoal 2 is:
\[ \neg (\exists x : D, \neg P x) \rightarrow \exists x : D, P x \rightarrow \forall x0 : D, P x0 \]

We conclude in that case by considering Tom, since his drinking leads to a contradiction:

\[ \text{Coq < exists Tom; intro Tom\_drinks.} \]
\[ 2 \text{ subgoals} \]

\[ D : Set \]
\[ R : D \rightarrow D \rightarrow Prop \]
\[ P : D \rightarrow Prop \]
\[ d : D \]
\[ EM : \forall A : Prop, A \lor \neg A \]
\[ \text{Non\_drinker : \exists x : D, \neg P x} \]
\[ \text{Tom : D} \]
\[ \text{Tom\_does\_not\_drink : \neg P Tom} \]
\[ \text{Tom\_drinks : P Tom} \]

\[ \forall x : D, P x \]

subgoal 2 is:
\[ \neg (\exists x : D, \neg P x) \rightarrow \exists x : D, P x \rightarrow \forall x0 : D, P x0 \]

There are several ways in which we may eliminate a contradictory case; a simple one is to use the \textit{absurd} tactic as follows:

\[ \text{Coq < absurd (P Tom); trivial.} \]
\[ 1 \text{ subgoal} \]

\[ D : Set \]
\[ R : D \rightarrow D \rightarrow Prop \]
\[ P : D \rightarrow Prop \]
\[ d : D \]
\[ EM : \forall A : Prop, A \lor \neg A \]

\[ \neg (\exists x : D, \neg P x) \rightarrow \exists x : D, P x \rightarrow \forall x0 : D, P x0 \]

We now proceed with the second case, in which actually any person will do; such a John Doe is given by the non-emptiness witness \( d \):
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Coq < intro No_nondrinker; exists d; intro d_drinks.
1 subgoal

D : Set
R : D -> D -> Prop
P : D -> Prop
d : D
EM : forall A : Prop, A \or ~ A
No_nondrinker : ~ (exists x : D, ~ P x)
d_drinks : P d

Now we consider any Dick in the bar, and reason by cases according to its drinking or not:

Coq < intro Dick; elim (EM (P Dick)); trivial.
1 subgoal

D : Set
R : D -> D -> Prop
P : D -> Prop
d : D
EM : forall A : Prop, A \or ~ A
No_nondrinker : ~ (exists x : D, ~ P x)
d_drinks : P d
Dick : D

The only non-trivial case is again treated by contradiction:

Coq < intro Dick_does_not_drink; absurd (exists x, ~ P x); trivial.
1 subgoal

D : Set
R : D -> D -> Prop
P : D -> Prop
d : D
EM : forall A : Prop, A \or ~ A
No_nondrinker : ~ (exists x : D, ~ P x)
d_drinks : P d
Dick : D
Dick_does_not_drink : ~ P Dick
exists x : D, ~ P x

Coq < exists Dick; trivial.
No more subgoals.

Coq < Qed.
(elim (EM (exists x, ~ P x))).
(intro Non_drinker; elim Non_drinker;
  intros Tom Tom_does_not_drink).
(exists Tom; intro Tom_drinks).
(absurd (P Tom); trivial).
(intro No_nondrinker; exists d; intro d_drinks).
(intro Dick; elim (EM (P Dick)); trivial).
(intro Dick_does_not_drink; absurd (exists x, ~ P x);
  trivial).
(exists Dick; trivial).
Qed.

Now, let us close the main section and look at the complete statements we
proved:

Coq < End Predicate_calculus.

Coq < Check refl_if.
refl_if :
  forall (D : Set) (R : D -> D -> Prop),
  (forall x y : D, R x y -> R y x) ->
  (forall x y z : D, R x y -> R y z -> R x z) ->
  forall x : D, (exists y : D, R x y) -> R x x

Coq < Check weird.
weird :
  forall (D : Set) (P : D -> Prop),
  D -> (forall x : D, P x) -> exists a : D, P a

Coq < Check drinker.
drinker :
  forall (D : Set) (P : D -> Prop),
  D ->
  (forall A : Prop, A \^\ A) ->
  exists x : D, P x -> forall x0 : D, P x0

Remark how the three theorems are completely generic in the most general
fashion; the domain D is discharged in all of them, R is discharged in refl_if only,
P is discharged only in weird and drinker, along with the hypothesis that D is
inhabited. Finally, the excluded middle hypothesis is discharged only in drinker.

Note that the name d has vanished as well from the statements of weird and
drinker, since Coq’s pretty-printer replaces systematically a quantification such
as forall d:D, E, where d does not occur in E, by the functional notation D->E.  
Similarly the name EM does not appear in drinker.

Actually, universal quantification, implication, as well as function formation,
are all special cases of one general construct of type theory called dependent prod-
uct. This is the mathematical construction corresponding to an indexed family of
functions. A function \( f \in \Pi x : D \cdot Cx \) maps an element x of its domain D to its (in-
dexed) codomain Cx. Thus a proof of \( \forall x : D \cdot Px \) is a function mapping an element
x of D to a proof of proposition Px.
1.4. Flexible use of local assumptions

Very often during the course of a proof we want to retrieve a local assumption and reintroduce it explicitly in the goal, for instance in order to get a more general induction hypothesis. The tactic generalize is what is needed here:

Coq < Section Predicate_Calculus.
Coq < Variables P Q : nat -> Prop.
P is declared
Q is declared
Coq < Variable R : nat -> nat -> Prop.
R is declared
Coq < Lemma PQR :
  forall x y:nat, (R x x -> P x -> Q x) -> P x -> R x y -> Q x.
1 subgoal

P, Q : nat -> Prop
R : nat -> nat -> Prop
floated forall x y : nat, (R x x -> P x -> Q x) -> P x -> R x y -> Q x.
1 subgoal

Coq < intros.
1 subgoal

P, Q : nat -> Prop
R : nat -> nat -> Prop
x, y : nat
H : R x x -> P x -> Q x
H0 : P x
H1 : R x y
floated Q x.
1 subgoal

Coq < generalize H0.
1 subgoal

P, Q : nat -> Prop
R : nat -> nat -> Prop
x, y : nat
H : R x x -> P x -> Q x
H0 : P x
H1 : R x y
floated P x -> Q x.

Sometimes it may be convenient to use a lemma, although we do not have a direct way to appeal to such an already proven fact. The tactic cut permits to use the lemma at this point, keeping the corresponding proof obligation as a new subgoal:
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Coq < cut (R x x); trivial.
1 subgoal

\[ P, Q : \text{n} \rightarrow \text{Prop} \]
\[ R : \text{n} \rightarrow \text{n} \rightarrow \text{Prop} \]
\[ x, y : \text{n} \]
\[ H : R \ x \ x \rightarrow P \ x \rightarrow Q \ x \]
\[ H0 : P \ x \]
\[ H1 : R \ x \ y \]

We clean the goal by doing an \texttt{Abort} command.

Coq < Abort.

1.4.5 Equality

The basic equality provided in COQ is Leibniz equality, noted infix like \( x = y \), when \( x \) and \( y \) are two expressions of type the same \textit{Set}. The replacement of \( x \) by \( y \) in any term is effected by a variety of tactics, such as \texttt{rewrite} and \texttt{replace}.

Let us give a few examples of equality replacement. Let us assume that some arithmetic function \( f \) is null in zero:

Coq < Variable \( f : \text{n} \rightarrow \text{n} \).
\( f \) is declared

Coq < Hypothesis \( \text{foo} : f \ 0 = 0 \).
\( \text{foo} \) is declared

We want to prove the following conditional equality:

Coq < Lemma \( L1 : \forall k : \text{n}, k = 0 \rightarrow f \ k = k \).

As usual, we first get rid of local assumptions with \texttt{intro}:

Coq < intros \( k \) \( E \).
1 subgoal

\[ P, Q : \text{n} \rightarrow \text{Prop} \]
\[ R : \text{n} \rightarrow \text{n} \rightarrow \text{Prop} \]
\[ f : \text{n} \rightarrow \text{n} \]
\[ \text{foo} : f \ 0 = 0 \]
\[ k : \text{n} \]
\[ E : k = 0 \]

Let us now use equation \( E \) as a left-to-right rewriting:
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Coq < rewrite E.
1 subgoal

P, Q : nat -> Prop
R : nat -> nat -> Prop
f : nat -> nat
foo : f 0 = 0
k : nat
E : k = 0

This replaced both occurrences of k by 0.

Now apply foo will finish the proof:

Coq < apply foo.
No more subgoals.

Coq < Qed.
(intros k E).
(rewrite E).
(apply foo).
Qed.
L1 is defined

When one wants to rewrite an equality in a right to left fashion, we should use rewrite <- E rather than rewrite E or the equivalent rewrite -> E. Let us now illustrate the tactic replace.

Coq < Hypothesis f10 : f 1 = f 0.
f10 is declared

Coq < Lemma L2 : f (f 1) = 0.
1 subgoal

P, Q : nat -> Prop
R : nat -> nat -> Prop
f : nat -> nat
foo : f 0 = 0
f10 : f 1 = f 0

Coq < replace (f 1) with 0.
2 subgoals
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\[ f \ 0 = 0 \]
\[ \text{subgoal 2 is:} \]
\[ 0 = f \ 1 \]

What happened here is that the replacement left the first subgoal to be proved, but another proof obligation was generated by the \texttt{replace} tactic, as the second subgoal. The first subgoal is solved immediately by applying lemma \texttt{foo}; the second one transitivity and then symmetry of equality, for instance with tactics \texttt{transitivity} and \texttt{symmetry}:

\begin{verbatim}
Coq < apply foo.
1 subgoal

\[ P, Q : \text{nat} \rightarrow \text{Prop} \]
\[ R : \text{nat} \rightarrow \text{nat} \rightarrow \text{Prop} \]
\[ f : \text{nat} \rightarrow \text{nat} \]
\[ \text{foo} : f \ 0 = 0 \]
\[ f10 : f \ 1 = f \ 0 \]

\[ ================ \]
\[ 0 = f \ 1 \]
\end{verbatim}

\begin{verbatim}
Coq < transitivity (f \ 0); symmetry; trivial.
No more subgoals.
\end{verbatim}

In case the equality \( t = u \) generated by \texttt{replace u with t} is an assumption (possibly modulo symmetry), it will be automatically proved and the corresponding goal will not appear. For instance:

\begin{verbatim}
Coq < Restart.
1 subgoal

\[ P, Q : \text{nat} \rightarrow \text{Prop} \]
\[ R : \text{nat} \rightarrow \text{nat} \rightarrow \text{Prop} \]
\[ f : \text{nat} \rightarrow \text{nat} \]
\[ \text{foo} : f \ 0 = 0 \]
\[ f10 : f \ 1 = f \ 0 \]

\[ ================ \]
\[ f (f \ 1) = 0 \]
\end{verbatim}

\begin{verbatim}
Coq < replace (f \ 0) with 0.
1 subgoal

\[ P, Q : \text{nat} \rightarrow \text{Prop} \]
\[ R : \text{nat} \rightarrow \text{nat} \rightarrow \text{Prop} \]
\[ f : \text{nat} \rightarrow \text{nat} \]
\[ \text{foo} : f \ 0 = 0 \]
\[ f10 : f \ 1 = f \ 0 \]

\[ ================ \]
\[ f (f \ 1) = 0 \]
\end{verbatim}

\begin{verbatim}
Coq < rewrite f10; rewrite foo; trivial.
No more subgoals.
\end{verbatim}
1.5. USING DEFINITIONS

The development of mathematics does not simply proceed by logical argumentation from first principles: definitions are used in an essential way. A formal development proceeds by a dual process of abstraction, where one proves abstract statements in predicate calculus, and use of definitions, which in the contrary one instantiates general statements with particular notions in order to use the structure of mathematical values for the proof of more specialised properties.

1.5.1 Unfolding definitions

Assume that we want to develop the theory of sets represented as characteristic predicates over some universe $U$. For instance:

Coq < Variable U : Type.
U is declared
Coq < Definition set := U -> Prop.
set is defined
Coq < Definition element (x:U) (S:set) := S x.
element is defined
Coq < Definition subset (A B:set) :=
   forall x:U, element x A -> element x B.
subset is defined

Now, assume that we have loaded a module of general properties about relations over some abstract type $T$, such as transitivity:

Coq < Definition transitive (T:Type) (R:T -> T -> Prop) :=
   forall x y z:T, R x y -> R y z -> R x z.
transitive is defined

Now, assume that we want to prove that subset is a transitive relation.

Coq < Lemma subset_transitive : transitive set subset.
1 subgoal

$P, Q : \text{nat -> Prop}$
$R : \text{nat -> nat -> Prop}$
$f : \text{nat -> nat}$
$\text{foo} : f 0 = 0$
$f 10 : f 1 = f 0$
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\[
U : \text{Type} \\
\text{-------------------} \\
\text{transitive set subset}
\]

In order to make any progress, one needs to use the definition of *transitive*. The `unfold` tactic, which replaces all occurrences of a defined notion by its definition in the current goal, may be used here.

Coq < `unfold transitive.`
1 subgoal

\[
P, Q : \text{nat -> Prop} \\
R : \text{nat -> nat -> Prop} \\
f : \text{nat -> nat} \\
foo : f 0 = 0 \\
f10 : f 1 = f 0 \\
U : \text{Type} \\
\text{-------------------} \\
\forall x y z : \text{set}, \\
\text{subset x y -> subset y z -> subset x z}
\]

Now, we must unfold `subset`:

Coq < `unfold subset.`
1 subgoal

\[
P, Q : \text{nat -> Prop} \\
R : \text{nat -> nat -> Prop} \\
f : \text{nat -> nat} \\
foo : f 0 = 0 \\
f10 : f 1 = f 0 \\
U : \text{Type} \\
\text{-------------------} \\
\forall x y z : \text{set}, \\
(\forall x0 : U, \text{element x0 x -> element x0 y}) \rightarrow \\
(\forall x0 : U, \text{element x0 y -> element x0 z}) \rightarrow \\
(\forall x0 : U, \text{element x0 x -> element x0 z})
\]

Now, unfolding `element` would be a mistake, because indeed a simple proof can be found by `auto`, keeping `element` an abstract predicate:

Coq < `auto.`
No more subgoals.

Many variations on `unfold` are provided in Coq. For instance, we may selectively unfold one designated occurrence:

Coq < `Undo 2.`
1 subgoal

\[
P, Q : \text{nat -> Prop} \\
R : \text{nat -> nat -> Prop}
\]
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\[ f : \text{nat} \rightarrow \text{nat} \]
\[ \text{foo} : f \ 0 = 0 \]
\[ f10 : f \ 1 = f \ 0 \]
\[ U : \text{Type} \]

\[ =\dec {\forall x \ y \ z : \text{set}, \quad \text{subset} \ x \ y \to \text{subset} \ y \ z \to \text{subset} \ x \ z} \]

One may also unfold a definition in a given local hypothesis, using the \texttt{in} notation:

\[ \text{Coq < unfold subset in H.} \]
1 subgoal

\[ P, \ Q : \text{nat} \rightarrow \text{Prop} \]
\[ R : \text{nat} \rightarrow \text{nat} \rightarrow \text{Prop} \]
\[ f : \text{nat} \rightarrow \text{nat} \]
\[ \text{foo} : f \ 0 = 0 \]
\[ f10 : f \ 1 = f \ 0 \]
\[ U : \text{Type} \]
\[ x, \ y, \ z : \text{set} \]
\[ \text{H} : \text{subset} \ x \ y \]
\[ \text{H0} : \forall x : U, \text{element} \ x \ y \rightarrow \text{element} \ x \ z \]

\[ =\dec {\text{subset} \ x \ z} \]

\[ \text{Coq < unfold subset in H.} \]
1 subgoal
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\[ H : \forall x0 : U, \text{element } x0 x \rightarrow \text{element } x0 y \]
\[ H0 : \forall x : U, \text{element } x y \rightarrow \text{element } x z \]
\[ \text{subset } x z \]

Finally, the tactic \texttt{red} does only unfolding of the head occurrence of the current goal:

Coq < \texttt{red}.
1 subgoal

\[ P, Q : \text{nat} \rightarrow \text{Prop} \]
\[ R : \text{nat} \rightarrow \text{nat} \rightarrow \text{Prop} \]
\[ f : \text{nat} \rightarrow \text{nat} \]
\[ \text{foo} : f \ 0 = 0 \]
\[ f10 : f \ 1 = f \ 0 \]
\[ U : \text{Type} \]
\[ x, y, z : \text{set} \]
\[ \text{H} : \forall x0 : U, \text{element } x0 x \rightarrow \text{element } x0 y \]
\[ \text{H0} : \forall x : U, \text{element } x y \rightarrow \text{element } x z \]
\[ \text{subset } x z \]

Coq < \texttt{auto}.
No more subgoals.

Coq < \texttt{Qed}.
(\texttt{unfold transitive}).
(\texttt{unfold subset at 2}).
(\texttt{intros **}).
(\texttt{unfold subset in H}).
(\texttt{red}).
\texttt{auto}.
\texttt{Qed}.
\texttt{subset\_transitive is defined}

1.5.2 Principle of proof irrelevance

Even though in principle the proof term associated with a verified lemma corre-
sponds to a defined value of the corresponding specification, such definitions can-
ot be unfolded in \texttt{COQ}: a lemma is considered an \texttt{opaque} definition. This con-
forms to the mathematical tradition of \textit{proof irrelevance}: the proof of a logical
proposition does not matter, and the mathematical justification of a logical devel-
opment relies only on \textit{provability} of the lemmas used in the formal proof.

Conversely, ordinary mathematical definitions can be unfolded at will, they are
\textit{transparent}. 
Chapter 2

Induction

2.1 Data Types as Inductively Defined Mathematical Collections

All the notions which were studied until now pertain to traditional mathematical logic. Specifications of objects were abstract properties used in reasoning more or less constructively; we are now entering the realm of inductive types, which specify the existence of concrete mathematical constructions.

2.1.1 Booleans

Let us start with the collection of booleans, as they are specified in the CoQ’s Prelude module:

Coq < Inductive bool : Set := true | false.
bool is defined
bool_rect is defined
bool_ind is defined
bool_rec is defined

Such a declaration defines several objects at once. First, a new Set is declared, with name bool. Then the constructors of this Set are declared, called true and false. Those are analogous to introduction rules of the new Set bool. Finally, a specific elimination rule for bool is now available, which permits to reason by cases on bool values. Three instances are indeed defined as new combinators in the global context: bool_ind, a proof combinator corresponding to reasoning by cases, bool_rec, an if-then-else programming construct, and bool_rect, a similar combinator at the level of types. Indeed:

Coq < Check bool_ind.
bool_ind
  : forall P : bool -> Prop,
    P true -> P false -> forall b : bool, P b

Coq < Check bool_rec.
bool_rec
CHAPTER 2. INDUCTION

: forall P : bool -> Set,
P true -> P false -> forall b : bool, P b

Coq < Check bool_rect.
bool_rect
: forall P : bool -> Type,
P true -> P false -> forall b : bool, P b

Let us for instance prove that every Boolean is true or false.

Coq < Lemma duality : forall b:bool, b = true / b = false.
1 subgoal

============================
forall b : bool, b = true / b = false

Coq < intro b.
1 subgoal

b : bool

b = true / b = false

We use the knowledge that b is a bool by calling tactic elim, which is this case will appeal to combinator bool_ind in order to split the proof according to the two cases:

Coq < elim b.
2 subgoals

b : bool

true = true / true = false
subgoal 2 is:
false = true / false = false

It is easy to conclude in each case:

Coq < left; trivial.
1 subgoal

b : bool

false = true / false = false

Coq < right; trivial.
No more subgoals.

Indeed, the whole proof can be done with the combination of the simple induction, which combines intro and elim, with good old auto:

Coq < Restart.
1 subgoal
2.1. DATA TYPES AS INDUCTIVELY DEFINED MATHEMATICAL COLLECTIONS

\\begin{align*}
&\forall b : \text{bool}, \ b = \text{true} \lor b = \text{false} \\
\text{Coq} &< \text{simple induction } b; \text{auto.} \\
\text{No more subgoals.} \\
\text{Coq} &< \text{Qed.} \\
&\text{(simple induction } b; \text{auto).} \\
\text{Qed.} \\
&\text{duality is defined}
\end{align*}

2.1.2 Natural numbers

Similarly to Booleans, natural numbers are defined in the Prelude module with constructors S and O:

\text{Coq} < \text{Inductive nat : Set :=} \\
&| O : \text{nat} \\
&| S : \text{nat -> nat.} \\
\text{nat is defined} \\
\text{nat_rect is defined} \\
\text{nat_ind is defined} \\
\text{nat_rec is defined}

The elimination principles which are automatically generated are Peano’s induction principle, and a recursion operator:

\text{Coq} < \text{Check nat_ind.} \\
\text{nat_ind} \\
&: \forall P : \text{nat -> Prop,} \\
&P O \rightarrow \\
&(\forall n : \text{nat, } P n \rightarrow P (S n)) \rightarrow \\
&\forall n : \text{nat, } P n \\
\text{Coq} < \text{Check nat_rec.} \\
\text{nat_rec} \\
&: \forall P : \text{nat -> Set,} \\
&P O \rightarrow \\
&(\forall n : \text{nat, } P n \rightarrow P (S n)) \rightarrow \\
&\forall n : \text{nat, } P n

Let us start by showing how to program the standard primitive recursion operator \text{prim_rec} from the more general \text{nat_rec}:

\text{Coq} < \text{Definition prim_rec := nat_rec (fun i:nat => nat).} \\
\text{prim_rec is defined} \\

That is, instead of computing for natural \text{i} an element of the indexed Set \text{(P i)}, \text{prim_rec} computes uniformly an element of \text{nat}. Let us check the type of \text{prim_rec}:
Chapter 2. Induction

Coq < About prim_rec.

\[ \text{prim_rec :} \]
\[ (\text{fun} \_ : \text{nat} \Rightarrow \text{nat}) \, \text{O} \rightarrow \]
\[ (\text{forall} \, n : \text{nat}, \]
\[ (\text{fun} \_ : \text{nat} \Rightarrow \text{nat}) \, n \rightarrow (\text{fun} \_ : \text{nat} \Rightarrow \text{nat}) \, (\text{S} \, n)) \rightarrow \]
\[ \text{forall} \, n : \text{nat}, (\text{fun} \_ : \text{nat} \Rightarrow \text{nat}) \, n \]

Argument scopes are \([\_ \text{ function scope } \_]\)

\[ \text{prim_rec is transparent} \]

Expands to: Constant Top.prim_rec

Oops! Instead of the expected type \(\text{nat} \rightarrow (\text{nat} \rightarrow \text{nat}) \rightarrow \text{nat} \rightarrow \text{nat}\) we get an apparently more complicated expression. Indeed the type of \(\text{prim_rec}\) is equivalent by rule \(\beta\) to its expected type; this may be checked in COQ by command \text{Eval Cbv Beta}, which \(\beta\)-reduces an expression to its normal form:

Coq < Eval cbv beta in

\[ ((\text{fun} \_ : \text{nat} \Rightarrow \text{nat}) \, \text{O} \rightarrow \]
\[ (\text{forall} \, y : \text{nat}, \]
\[ (\text{fun} \_ : \text{nat} \Rightarrow \text{nat}) \, y \rightarrow (\text{fun} \_ : \text{nat} \Rightarrow \text{nat}) \, (\text{S} \, y)) \rightarrow \]
\[ \text{forall} \, n : \text{nat}, (\text{fun} \_ : \text{nat} \Rightarrow \text{nat}) \, n \].

\[ = \text{nat} \rightarrow (\text{nat} \rightarrow \text{nat} \rightarrow \text{nat}) \rightarrow \text{nat} \rightarrow \text{nat} \]

: \text{Set}

Let us now show how to program addition with primitive recursion:

Coq < Definition addition (n m : nat) :=
\[ \text{prim_rec m} \, (\text{fun p rec : nat} \Rightarrow \text{S} \, \text{rec}) \, n. \]

addition is defined

That is, we specify that \((\text{addition} \, n \, m)\) computes by cases on \(n\) according to its main constructor; when \(n = \text{O}\), we get \(m\); when \(n = \text{S} \, p\), we get \((\text{S} \, \text{rec})\), where \(\text{rec}\) is the result of the recursive computation \((\text{addition} \, p \, m)\). Let us verify it by asking COQ to compute for us say \(2 + 3\):

Coq < Eval compute in

\[ (\text{addition} \, (\text{S} \, (\text{S} \, \text{O}))) \, (\text{S} \, \text{O})). \]

\[ = \text{S} \, (\text{S} \, \text{S} \, (\text{S} \, \text{S} \, \text{O})). \]

\[ : (\text{fun} \_ : \text{nat} \Rightarrow \text{nat}) \, (\text{S} \, \text{S} \, \text{O}). \]

Actually, we do not have to do all explicitly. COQ provides a special syntax \text{Fixpoint/match} for generic primitive recursion, and we could thus have defined directly addition as:

Coq < Fixpoint plus (n m : nat) {struct n} : nat :=
\[ \text{match n with} \]
\[ | \text{O} \Rightarrow m \]
\[ | \text{S} \, p \Rightarrow \text{S} \, (\text{plus} \, p \, m) \]
\[ \text{end.} \]

plus is defined

plus is recursively defined (decreasing on 1st argument)
2.1. DATA TYPES AS INDUCTIVELY DEFINED MATHEMATICAL COLLECTIONS

For the rest of the session, we shall clean up what we did so far with types bool and nat, in order to use the initial definitions given in COQ’s Prelude module, and not to get confusing error messages due to our redefinitions. We thus revert to the initial state:

Coq < Reset Initial.

2.1.3 Simple proofs by induction

Let us now show how to do proofs by structural induction. We start with easy properties of the plus function we just defined. Let us first show that \( n = n + 0 \).

Coq < Lemma plus_n_O : forall n:nat, n = n + 0.
1 subgoal

\[ \forall n \in \text{nat}, \ n = n + 0 \]

Coq < intro n; elim n.
2 subgoals

\[ n \in \text{nat} \]

\[ 0 = 0 + 0 \]

subgoal 2 is:

\[ \forall n0 \in \text{nat}, n0 = n0 + 0 \rightarrow S n0 = S n0 + 0 \]

What happened was that elim \( n \), in order to construct a Prop (the initial goal) from a nat (i.e. \( n \)), appealed to the corresponding induction principle nat_ind which we saw was indeed exactly Peano’s induction scheme. Pattern-matching instantiated the corresponding predicate \( P \) to \( \text{fun} \ n \in \text{nat} \Rightarrow n = n \ 0+ \), and we get as subgoals the corresponding instantiations of the base case \( (P \ 0) \), and of the inductive step \( \forall y : \text{nat}, P y \rightarrow P (S y) \). In each case we get an instance of function plus in which its second argument starts with a constructor, and is thus amenable to simplification by primitive recursion. The COQ tactic simpl can be used for this purpose:

Coq < simpl.
2 subgoals

\[ n \in \text{nat} \]

\[ 0 = 0 \]

subgoal 2 is:

\[ \forall n0 : \text{nat}, n0 = n0 + 0 \rightarrow S n0 = S n0 + 0 \]

Coq < auto.
1 subgoal

\[ n \in \text{nat} \]

\[ \forall n0 : \text{nat}, n0 = n0 + 0 \rightarrow S n0 = S n0 + 0 \]
We proceed in the same way for the base step:

Coq < simpl; auto.
No more subgoals.
Coq < Qed.
(intro n; elim n).
(simpl1).
auto.
(simpl; auto).
Qed.

plus_n_O is defined

Here auto succeeded, because it used as a hint lemma eq_S, which say that successor preserves equality:

Coq < Check eq_S.
eq_S
: forall x y : nat, x = y -> S x = S y

Actually, let us see how to declare our lemma plus_n_O as a hint to be used by auto:

Coq < Hint Resolve plus_n_O .

We now proceed to the similar property concerning the other constructor S:

Coq < Lemma plus_n_S : forall n m:nat, S (n + m) = n + S m.
1 subgoal

forall n m : nat, S (n + m) = n + S m

We now go faster, remembering that tactic simple induction does the necessary intros before applying elim. Factoring simplification and automation in both cases thanks to tactic composition, we prove this lemma in one line:

Coq < simple induction n; simpl; auto.
No more subgoals.
Coq < Qed.
(simple induction n; simpl; auto).
Qed.

plus_n_S is defined

Coq < Hint Resolve plus_n_S .

Let us end this exercise with the commutativity of plus:

Coq < Lemma plus_com : forall n m:nat, n + m = m + n.
1 subgoal

forall n m : nat, n + m = m + n
2.1. DATA TYPES AS INDUCTIVELY DEFINED MATHEMATICAL COLLECTIONS

Here we have a choice on doing an induction on \( n \) or on \( m \), the situation being symmetric. For instance:

Coq < simple induction m; simpl; auto.
1 subgoal

\( n, m : \text{nat} \)

\[ \forall n0 : \text{nat}, n + n0 = n0 + n \rightarrow n + S n0 = S (n0 + n) \]

Here \text{auto} succeeded on the base case, thanks to our hint \text{plus_n_O}, but the induction step requires rewriting, which \text{auto} does not handle:

Coq < intros m' E; rewrite <- E; auto.
No more subgoals.

Coq < Qed.
(simple induction m; simpl; auto).
[intros m' E; rewrite <- E; auto].
Qed.
\text{plus_com} is defined

2.1.4 Discriminate

It is also possible to define new propositions by primitive recursion. Let us for instance define the predicate which discriminates between the constructors \( \text{O} \) and \( \text{S} \): it computes to \text{False} when its argument is \( \text{O} \), and to \text{True} when its argument is of the form \( (S n) \):

Coq < Definition Is_S (n:nat) := match n with
| O => False
| S p => True
end.
\text{Is_S} is defined

Now we may use the computational power of \text{Is_S} in order to prove trivially that \( (Is_S (S n)) \):

Coq < Lemma S_Is_S : forall n:nat, Is_S (S n).
1 subgoal

\[ \forall n : \text{nat}, Is_S (S n) \]

Coq < simpl; trivial.
No more subgoals.

Coq < Qed.
(simpl; trivial).
Qed.
\text{S_Is_S} is defined
But we may also use it to transform a \texttt{False} goal into \texttt{(Is\_S 0)}. Let us show a particularly important use of this feature; we want to prove that 0 and \texttt{S} construct different values, one of Peano’s axioms:

\[
\text{Coq < Lemma no\_confusion : forall n:nat, 0 <> S n.}
\]

\[
1 \text{ subgoal}
\]

\[
\begin{array}{l}
\text{forall n : nat, 0 <> S n} \\
\end{array}
\]

First of all, we replace negation by its definition, by reducing the goal with tactic \texttt{red}; then we get contradiction by successive \texttt{intros}:

\[
\text{Coq < red; intros n H.}
\]

\[
1 \text{ subgoal}
\]

\[
\begin{array}{l}
n : \text{nat} \\
H : 0 = S n \\
\end{array}
\]

\[
\begin{array}{l}
\text{False} \\
\end{array}
\]

Now we use our trick:

\[
\text{Coq < change (Is\_S 0).}
\]

\[
1 \text{ subgoal}
\]

\[
\begin{array}{l}
n : \text{nat} \\
H : 0 = S n \\
\end{array}
\]

\[
\text{Is\_S 0}
\]

Now we use equality in order to get a subgoal which computes out to \texttt{True}, which finishes the proof:

\[
\text{Coq < rewrite H; trivial.}
\]

\[
1 \text{ subgoal}
\]

\[
\begin{array}{l}
n : \text{nat} \\
H : 0 = S n \\
\end{array}
\]

\[
\text{Is\_S (S n)}
\]

\[
\text{Coq < simpl; trivial.}
\]

\[
\text{No more subgoals.}
\]

Actually, a specific tactic \texttt{discriminate} is provided to produce mechanically such proofs, without the need for the user to define explicitly the relevant discrimination predicates:

\[
\text{Coq < Restart.}
\]

\[
1 \text{ subgoal}
\]
2.2. LOGIC PROGRAMMING

forall n : nat, 0 <> S n
Coq < intro n; discriminate.
No more subgoals.
Coq < Qed.
(intro n; discriminate).
Qed.
no_confusion is defined

2.2 Logic programming

In the same way as we defined standard data-types above, we may define inductive
families, and for instance inductive predicates. Here is the definition of predicate
≤ over type nat, as given in COQ’s Prelude module:

Coq < Inductive le (n:nat) : nat -> Prop :=
  | le_n : le n n
  | le_S : forall m:nat, le n m -> le n (S m).

This definition introduces a new predicate le:nat->nat->Prop, and the two
constructors le_n and le_S, which are the defining clauses of le. That is, we
get not only the “axioms” le_n and le_S, but also the converse property, that
(le n m) if and only if this statement can be obtained as a consequence of these
defining clauses; that is, le is the minimal predicate verifying clauses le_n and
le_S. This is insured, as in the case of inductive data types, by an elimination prin-
ciple, which here amounts to an induction principle le_ind, stating this minimality
property:

Coq < Check le.
le :
  nat -> nat -> Prop
Coq < Check le_ind.
le_ind :
  forall (n : nat) (P : nat -> Prop),
  P n ->
  (forall m : nat, le n m -> P m -> P (S m)) ->
  forall n0 : nat, le n n0 -> P n0

Let us show how proofs may be conducted with this principle. First we show
that \( n \leq m \Rightarrow n + 1 \leq m + 1 \):

Coq < Lemma le_n_S : forall n m:nat, le n m -> le (S n) (S m).
1 subgoal

forall n m : nat, le n m -> le (S n) (S m)
Coq < intros n m n_le_m.
1 subgoal
\( n, m : \text{nat} \\
\text{n_le_m : le n m} \\
\text{============================} \\
\text{le (S n) (S m)} \)

Coq \(<\) elim n_le_m.
2 subgoals

\( n, m : \text{nat} \\
\text{n_le_m : le n m} \\
\text{============================} \\
\text{le (S n) (S n)} \\
\text{subgoal 2 is:} \\
\text{forall m0 : nat,} \\
\text{le n m0 -> le (S n) (S m0) -> le (S n) (S (S m0))} \)

What happens here is similar to the behaviour of \text{elim} on natural numbers: it appeals to the relevant induction principle, here \text{le_ind}, which generates the two subgoals, which may then be solved easily with the help of the defining clauses of \text{le}.

Coq \(<\) apply le_n; trivial.
1 subgoal

\( n, m : \text{nat} \\
\text{n_le_m : le n m} \\
\text{============================} \\
\text{forall m0 : nat,} \\
\text{le n m0 -> le (S n) (S m0) -> le (S n) (S (S m0))} \)

Coq \(<\) intros; apply le_S; trivial.
No more subgoals.

Now we know that it is a good idea to give the defining clauses as hints, so that the proof may proceed with a simple combination of \text{induction and auto}.

Coq \(<\) Restart.
I subgoal

\( \text{=================================} \) \\
\text{forall n m : nat, le n m -> le (S n) (S m)} \\
Coq \(<\) Hint Resolve le_n le_S .

We have a slight problem however. We want to say “Do an induction on hypothesis \((\text{le n m})\)”, but we have no explicit name for it. What we do in this case is to say “Do an induction on the first unnamed hypothesis”, as follows.

Coq \(<\) simple induction 1; auto.
No more subgoals.
Coq \(<\) Qed.
2.2. LOGIC PROGRAMMING

Hint Resolve le_n le_S.
(simple induction 1; auto).
Qed.
le_n_S is defined

Here is a more tricky problem. Assume we want to show that \( n \leq 0 \Rightarrow n = 0 \).
This reasoning ought to follow simply from the fact that only the first defining
clause of le applies.

Coq < Lemma tricky : forall n:nat, le n 0 -> n = 0.
1 subgoal

============
forall n : nat, le n 0 -> n = 0

However, here trying something like induction 1 would lead nowhere (try it
and see what happens). An induction on \( n \) would not be convenient either. What we
must do here is analyse the definition of le in order to match hypothesis (le n 0)
with the defining clauses, to find that only \( le_n \) applies, whence the result. This
analysis may be performed by the “inversion” tactic inversion_clear as follows:

Coq < intros n H; inversion_clear H.
1 subgoal

n : nat
============
0 = 0

Coq < trivial.
No more subgoals.

Coq < Qed.
(intros n H; inversion_clear H).
trivial.
Qed.
tricky is defined
Chapter 3

Modules

3.1 Opening library modules

When you start Coq without further requirements in the command line, you get a bare system with few libraries loaded. As we saw, a standard prelude module provides the standard logic connectives, and a few arithmetic notions. If you want to load and open other modules from the library, you have to use the Require command, as we saw for classical logic above. For instance, if you want more arithmetic constructions, you should request:

Coq < Require Import Arith.

Such a command looks for a (compiled) module file Arith.vo in the libraries registered by Coq. Libraries inherit the structure of the file system of the operating system and are registered with the command Add LoadPath. Physical directories are mapped to logical directories. Especially the standard library of Coq is pre-registered as a library of name Coq. Modules have absolute unique names denoting their place in Coq libraries. An absolute name is a sequence of single identifiers separated by dots. E.g. the module Arith has full name Coq.Arith.Arith and because it resides in eponym subdirectory Arith of the standard library, it can be as well required by the command


This may be useful to avoid ambiguities if somewhere, in another branch of the libraries known by Coq, another module is also called Arith. Notice that by default, when a library is registered, all its contents, and all the contents of its subdirectories recursively are visible and accessible by a short (relative) name as Arith. Notice also that modules or definitions not explicitly registered in a library are put in a default library called Top.

The loading of a compiled file is quick, because the corresponding development is not type-checked again.
3.2 Creating your own modules

You may create your own module files, by writing COQ commands in a file, say my_module.v. Such a module may be simply loaded in the current context, with command Load my_module. It may also be compiled, in “batch” mode, using the UNIX command coqc. Compiling the module my_module.v creates a file my_module.vo that can be reloaded with command Require Import my_module.

If a required module depends on other modules then the latters are automatically required beforehand. However their contents is not automatically visible. If you want a module $M$ required in a module $N$ to be automatically visible when $N$ is required, you should use Require Export $M$ in your module $N$.

3.3 Managing the context

It is often difficult to remember the names of all lemmas and definitions available in the current context, especially if large libraries have been loaded. A convenient Search command is available to lookup all known facts concerning a given predicate. For instance, if you want to know all the known lemmas about both the successor and the less or equal relation, just ask:

```
Coq < Search S le.
```

```
le_S: forall n m : nat, n <= m -> n <= S m
le_ind:
  forall (n : nat) (P : nat -> Prop),
  P n ->
  (forall m : nat, n <= m -> P m -> P (S m)) ->
  forall n0 : nat, n <= n0 -> P n0
le_S_n: forall n m : nat, S n <= S m -> n <= m
le_n_S: forall n m : nat, n <= m -> S n <= S m
```

Another command SearchHead displays only lemmas where the searched predicate appears at the head position in the conclusion.

```
Coq < SearchHead le.
```

```
le_n: forall n : nat, n <= n
le_S: forall n m : nat, n <= m -> n <= S m
le_pred:
  forall n m : nat, n <= m -> Nat.pred n <= Nat.pred m
le_S_n: forall n m : nat, S n <= S m -> n <= m
le_0_n: forall n : nat, 0 <= n
le_n_S: forall n m : nat, n <= m -> S n <= S m
```

A new and more convenient search tool is SearchPattern developed by Yves Bertot. It allows finding the theorems with a conclusion matching a given pattern, where _ can be used in place of an arbitrary term. We remark in this example, that COQ provides usual infix notations for arithmetic operators.

```
Coq < SearchPattern (_ + _ = _).
```
3.4. NOW YOU ARE ON YOUR OWN

This tutorial is necessarily incomplete. If you wish to pursue serious proving in COQ, you should now get your hands on COQ’s Reference Manual, which contains a complete description of all the tactics we saw, plus many more. You also should look in the library of developed theories which is distributed with COQ, in order to acquaint yourself with various proof techniques.